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Two-Tier Cache-Aided Full-Duplex Hybrid Satellite–Terrestrial Communication Networks

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Abstract— Enabling global Internet access is challenging for cellular-based Internet of Things (IoT) due to the limited range of terrestrial network services. One viable solution is to deploy IoT over satellite systems for coverage extension. However, operating a hybrid satellite-terrestrial network (STN) might incur high satellite bandwidth consumption and excessive service latency. Aiming to reduce the content delivery latency from the Internet-connected gateway to the users, this work proposes a two-tier cache-enabled model with full-duplex transmissions where content caches are deployed at the satellite and ground station. A closed-form solution for the successful delivery probability (SDP) of the files is derived considering the requested content distributions and channel statistics. Then, the SDP performance under common caching policies can be evaluated. The results are also used to optimize cache placement under caching capacity constraints. Numerical results demonstrate the performance improvements of the proposed system over those of single-tier cache-aided and half-duplex transmission systems.

Index Terms— caching, full-duplex, half-duplex, satellite IoT, successful delivery probability

I. Introduction

ALTHOUGH being around for many years, cellular-based technologies for the Internet of Things (IoT) have not fulfilled the demand of globally connecting everyone and everything. To offer service across all geographic

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regions, integrating satellites into IoT networks has been touted. With better service reliability and coverage, satellite communications can provide more effective solutions for IoT than cellular networks. For instance, satellites can provide more than 99.9% availability with proper constellation arrangements [1] and can cover an enormous area of the Earth that is suitable for bringing connectivity to rural or remote areas. Reflecting the utilization of satellite systems as a promising solution complimenting terrestrial networks, many projects have been developed such as the new radio supporting integrated satellite–terrestrial networks by 3GPP [2], the Satellite and Terrestrial Network for 5G (SaT5G) [3], SpaceX’s Starlink satellite project [4].

Operating a hybrid satellite–terrestrial network (STN), however, poses challenges in service delays and satellite bandwidth costs. A content delivered to end-users from the Internet-connected gateway is relayed through satellite(s) and ground station(s), which extends the serving time in addition to very pricey and often limited satellite bandwidth. Excessive content delivery delay can hinder potential use of STN to support delay-sensitive applications. To address these challenges, caching has been proposed as a mechanism to relocate content storage to edge devices that are closer to users [5]–[12]. In wireless networks, edge caching has been shown to shorten delivery latency and reduce network congestion [13]. In addition, caching contents in satellites or terrestrial stations helps avoid re-transmissions of the same contents and thus significantly reducing in-network traffic [14].

In this paper¹, we study a two-tier cache-aided STN model considering end-to-end communications (i.e., from the Internet-connected gateway via the satellite and ground station to the end-users). Our work focuses on the content delivery latency analysis in terms of the successful delivery probability (SDP) with full-duplex (FD) transmissions at the satellite and ground station, albeit with imperfect self-interference (SI) cancellation. Deploying FD communications will potentially shorten the service delivery time compared to the half-duplex (HD) mode providing sufficiently effective SI cancellation, and hence increasing the SDP. The main contributions of this paper are summarized as follows:

1) A closed-form SDP expression for the two-tier cache-enabled STN is derived under the FD mode taking into account the requested content distributions, satellite/terrestrial realistic channel statistics, and caching configurations. Three commonly studied caching policies, namely uniform caching, ground station most popular caching, and satellite most popular caching, are also analyzed.

2) Based upon the derived results, the SDP performance of STN under different transmission modes (FD with and without FastForward relaying or HD mode) and

¹The preliminary results of this work [15] have been presented to the 10th International Conference on Computational Data and Social Networks.

single-tier cache-aided model can be investigated. The network capacity in terms of the maximum number of supportable users given a minimum SDP requirement is also studied.

3) An SDP maximization-based cache placement is designed subject to the caching capacity constraints of the two cache tiers. While the optimal solution is computationally intractable, low-complexity algorithms are developed to obtain caching policies with superior performance comparing to using common caching configurations.

4) Various simulation scenarios are considered to illustrate the improved SDP performance of the proposed system over its single-tier cache-aided counterparts and HD transmission systems under similar caching constraints and energy consumption.

The remainder of this paper is organized as follows. Section II introduces the related work on edge caching in STN. Section III describes the system model and the transmission scheme for FD transmissions. Section IV analyzes the SDP of the general caching and special caching policies, as well as several baseline schemes such as FD without FastForward relaying protocol, HD transmission scheme, and single-tier cache models. Cache placement design and numerical results are presented in Sections V and VI, respectively. Finally, conclusions are made in Section VII.

II. Edge caching techniques in STN

TWO types of satellites commonly used in STN are the low earth orbit (LEO) and geosynchronous equatorial orbit (GEO) satellites. Employing LEO satellites can achieve lower transmission latency than using GEO satellites. However, using LEO satellites poses a very challenging problem on dynamic network topology since they are only in contact with stations on the ground for a brief period due to their traveling speeds.

Taking into account the relative motion of LEO satellites and terrestrial users, [5] proposes a cross-time-slot graph for the time-varying topology file distribution in combination with the back-tracing partition directed on-path caching distribution mechanism (BPDPM) to reduce the caching overhead. The proposed BPDPM focuses on cache location decision using a modified Dijkstra routing algorithm and cache content placement using the popularity-aware multiple regions cooperative cache strategy. Proposing a time-evolving covering set in-network caching mechanism, [6] exploits the local connectivity of different regions in choosing the intermediate nodes with better connectivity that relay the objective file. The proposed mechanism includes a novel event updating graph that captures the time-varying topology of the network, a minimum time-evolving covering set algorithm for cache location decision, and a query-based minimum delay distribution path algorithm for optimal path finding during file access procedure. Leveraging the broadcast characteristic of satellite to efficiently provide traffic offloading for terrestrial networks, [7] proposes

an integrated LEO satellite/terrestrial cooperative caching scheme, in which the terrestrial access point is able to share its cached content with the satellite.

Employing GEO satellites in STN is suitable for latency-non-sensitive services due to the satellite high altitude. STN with GEO satellites have the advantage of the static network topology and non-interrupting transmission between the satellites and devices on the ground. Incorporating wireless caching into STN, [8] proposes an amplify-and-forward relaying protocols, where cache is enabled at relays on the ground. The proposal in [8] considers the most popular and uniform content-based cache placement schemes, which shows a substantial improvement in outage probability over the traditional approach without caching. To offload the backhaul of terrestrial networks, [9] proposes using hybrid STN in combination with an off-line edge caching algorithm. Caches are placed at base stations, in which each cache is divided into two parts. The first part is filled with the most local popular files whilst the other part is filled with the most global popular files. On the performance investigation of 5G edge caching over satellite, [10] uses off-line caching approaches to offload terrestrial backhaul networks. The proposal in [10] focuses on the multimode satellite backhaul in two use cases, i.e., dense urban areas and sparsely populated regions.

It should be emphasized that the system models in [5]–[10] are single-tier cache-enabled STN where caching is used for terrestrial devices only. Moreover, [8]–[10] consider dual-hop downlink satellite–relay(s)–users transmission only. The satellite uplink communications from the Internet-connected gateway to the satellite is not considered, which is usually the bottleneck in STN due to large bandwidth consumption.

Enabling cache capacity at satellite, [11], [12] propose two-tier caching models, where the first- and second-tier caches are placed in the terrestrial base stations and the satellite, respectively. The two-tier cache system will potentially: (i) reduce service latency since more contents are stored closer to the end-users; (ii) efficiently utilize satellite bandwidth as the contents are prefetched at the satellite during off-peak hours; Furthermore, when same contents are requested by multiple users, having the contents cached at satellite and broadcasting them at once can save satellite bandwidth. The work in [11] focuses on analyzing the satellite bandwidth consumption with an off-line caching scheme where popular contents are cached at each ground station for local services, while the satellite's cache is used for the most popular contents within its coverage (containing multiple ground stations) to take advantage of the satellite's broadcast nature to the ground stations. Non-cached contents can be retrieved from the gateway if needed. On the other hand, our current work focuses on investigating the file delivery time through the successful delivery probability (SDP) metric considering the realistic channel models. SDP analysis is typically relevant in delay-sensitive applications [16]. In addition, while [11] considers a scenario where a

file is fully cached, our work considers more general caching policies allowing partial caching of files. Hence, caching capability can be exploited better for improved performance. Employing the LEO satellite using a two-tier caching model, [12] designs a terrestrial link between the core network and base stations to maintain services when quality satellite connection is absent.

III. System model

CONSIDER a STN with a satellite S , an Internet-connected satellite gateway G , a ground station G_s , and a set of K ground end-users $U_i, i = 1, \dots, K$ as depicted in Fig. 1. In this STN, S is a GEO satellite and G_s is a low power ground base station equipped with a satellite receiver. Both S and G_s are cache enabled. Assuming there are no direct links from user U_i to satellite S and from the ground station to the gateway due to weather factors, long distance, and heavy shadowing.

Fig. 1: Satellite IoT networks architecture.

A. Two-tier Caching Model

Consider a two-tier caching model in which the first tier at G_s with storage capacity of C_1 (bits) and the second tier at S with storage capacity C_2 (bits). The gateway is connected to the Internet and hosts N files $W_1, \dots, W_N$, which are assumed to be of equal size of F (bits)². Without loss of generality, assume $C_i/F = N_i, i = 1, 2$ being integer numbers with $N_1 + N_2 < N$ as it is not possible to cache all files. The generic caching policy considered in this paper is represented by $\mu_1 = [\mu_{1,1}, \dots, \mu_{1,N}]$, $\mu_2 = [\mu_{2,1}, \dots, \mu_{2,N}]$ where $\mu_{1,n}, \mu_{2,n} \in [0, 1]$ are the portions of W_n fetched at the first-tier cache (at ground station) and the second-tier cache (at satellite), respectively. It is assumed that the

cached portions of the same file do not have overlapping contents. Other conditions of the caching policy are $\mu_{1,n} + \mu_{2,n} \leq 1, \forall n$ and $\sum_{n=1}^N \mu_{i,n} = N_i, i = 1, 2$ as it is optimal to utilize all caching capacities.

B. File Popularity Model

The request probability q_n for file W_n follows the Zipf distribution, i.e.,

$$q_n = \frac{n^{-\alpha}}{\sum_{m=1}^N m^{-\alpha}} \quad (1)$$

where $0 < \alpha < 1$ denotes the Zipf skewness factor [17]. A large value of α means the requests on high-popularity files, whereas a small value of α is related to the requests with heavy-tailed popularity. Without loss of generality, it is assumed that files $W_1, \dots, W_N$ have decreasing popularity.

C. Channel Model

We consider block-based communications, where a transmission section is accomplished within a coherence time T (seconds). Both large-scale fading and small-scale fading are considered. Large-scale fading is modeled by the distance-dependent power-law path-loss attenuation $d_{mn}^{-\alpha_i}$, where d_{mn} denotes the distance between nodes m and n , and α_i represents the path-loss exponent. For small-scale fading, the channel model proposed in [18] is commonly used for satellite terrestrial communications [8], [19]–[21], while Rayleigh fading channel is commonly used for terrestrial wireless communications.

The (uplink and downlink) satellite-terrestrial links have multipath fading and shadow fading. The shadow fading is composed of the line-of-sight (LOS) shadow fading and multiplicative shadow fading; The multipath fading composes of one LOS and many weak scatter components [18]. Therefore, the channel coefficient $\tilde{h}_S(t)$ for the satellite-terrestrial link can be modeled as

$$\tilde{h}_S(t) = A(t)e^{j\vartheta(t)} + Z(t)e^{j\zeta_0} \quad (2)$$

where $j^2 = -1$, $A(t)$ is the independent stationary random process representing the amplitude of the scatter components and following the Rayleigh distribution; $Z(t)$ represents the LOS component and follows the Nakagami distribution; $\vartheta(t)$ is the stationary random phase process with a uniform distribution over $[0, 2\pi)$; ζ_0 is the deterministic phase of the LOS component.

The normalized channel power gain $h_S(t) = |\tilde{h}_S(t)|^2/\sigma^2$, where σ^2 is the additive white Gaussian noise power, has the probability distribution function (PDF) $f_{h_S}(x)$ given by [18]

$$f_{h_S}(x) = \alpha_1 \exp\left(-\frac{x}{b_1}\right) {}_1F_1(m_1, 1, \sigma_1 x) \quad (3)$$

with coefficients $\alpha_1 = \left(\frac{b_1 m_1}{b_1 m_1 + \Omega_1}\right)^{m_1} \frac{1}{b_1}$, $\sigma_1 = \frac{\Omega_1}{b_1(b_1 m_1 + \Omega_1)}$, where $b_1 = \mathbb{E}[|A(t)|^2]$ represents the

²The proposed scheme can be readily extended for unequal file size setting.

average power of the scatter components ; $\Omega_1 = \mathbb{E}[|Z(t)|^2]$ represents the average power of the LOS component; $m_1 = \Omega_1^2 / \text{Var}[|Z(t)|^2]$ is the Nakagami parameter; ${}_1F_1(\cdot, \cdot, \cdot)$ is the confluent hypergeometric function of the first kind [22, eq.(9.210.1)]: ${}_1F_1(a; 1; z) = \sum_{k=0}^{a-1} \frac{(-1)^k (1-a)_k z^k e^z}{(k!)^2}$, where $(x)_n = \frac{\Gamma(x+n)}{\Gamma(x)}$ denotes the Pochhammer's symbol and $\Gamma(x)$ is Gamma function. Note that we omit the time-dependent index due to the stationarity assumption.

LEMMA 1. *The cumulative distribution function (CDF) $F_{h_S}(x)$ of $h_S(t)$ is*

$$F_{h_S}(x) = \sum_{k=0}^{m_1-1} \sum_{q=0}^k \frac{\alpha_1 e^{(\sigma_1-1/b_1)x} (-1)^{2k-q} (1-m_1)_k \sigma_1^k x^q}{k! q! (\sigma_1 - 1/b_1)^{k-q+1}}. \quad (4)$$

Proof:

$F_{h_S}(x)$ of $h_S(t)$ can be derived using elementary functions and finite summations

$$\begin{aligned} F_{h_S}(x) &= \int_0^x f_{h_S}(y) dy \\ &= \sum_{k=0}^{m_1-1} \frac{\alpha_1 (-1)^k (1-m_1)_k \sigma_1^k}{(k!)^2} \int_0^x y^k e^{(\sigma_1-1/b_1)y} dy. \end{aligned}$$

Using the following result on integration with exponential function:

$$\int y^k e^{cy} dy = e^{cy} \sum_{q=0}^k \frac{(-1)^{k-q} k!}{q! c^{k-q+1}} y^q,$$

we can derive (4). ■

To validate (4), Fig. 2 depicts the PDF and CDF of $h_S(t)$ with $\{m_1, b_1, \Omega_1\} = \{5, 0.251, 0.279\}$. We can see that as x becomes large, $F_{h_S}(x)$ approaches one as expected.

Fig. 2: PDF and CDF of $h_S(t)$.

The terrestrial link between the ground station and user U_i is modeled as a Rayleigh fading channel with the (normalized) channel power gain $h_i(t)$ governed by the

following distribution functions

$$f_{h_i}(x) = \frac{1}{\bar{h}_i} \exp\left(-\frac{x}{\bar{h}_i}\right), \quad F_{h_i}(x) = 1 - \exp\left(-\frac{x}{\bar{h}_i}\right) \quad (5)$$

with $\bar{h}_i$ is the average channel power gain.

D. Full-duplex Transmission with FastForward Relaying

Assume that the users will be served in a time-division multiple access (TDMA) manner which allows each user to be consecutively active in T/K (seconds). The time slot T/K will be large enough to ensure data freshness in practical IoT applications [23]. Under TDMA, inter-user interference does not exist. Both the satellite and ground station operate in the FD mode. When the channel state information is known at the transmitter, rate adaption is employed and the achievable rates (bps) on the links are

$$\begin{aligned} R_i &= B_{G_s} \log_2 \left(1 + P_{G_s} h_i d_i^{-\alpha_g} \right), \\ R_S &= B_S \log_2 \left(1 + \frac{P_S h_S d_S^{-\alpha_s}}{\beta P_{G_s} + 1} \right), \\ R_G &= B_G \log_2 \left(1 + \frac{P_G h_G d_G^{-\alpha_s}}{\beta P_S + 1} \right), \end{aligned} \quad (6)$$

where B_G , B_S , and B_{G_s} are the bandwidths (Hz) of the satellite uplink, downlink, and the terrestrial downlink, respectively; P_S , P_G , and P_{G_s} denote the transmit powers of the satellite, gateway, and ground station, respectively; d_i , d_S , and d_G are the distances of terrestrial link, satellite downlink and satellite uplink, respectively; and α_g , and α_s represent the path-loss exponents of the terrestrial and satellite links, respectively; βP_S and βP_{G_s} represent the normalized residual SI powers at the satellite and ground station respectively, which are assumed to be proportional to the transmit powers with coefficient $\beta \geq 0$ being the SI cancellation quality parameter [24]. For notational simplicity, denote:

$$\begin{aligned} \bar{P}_S &\triangleq P_S d_S^{-\alpha_s} / (\beta P_{G_s} + 1), \quad \bar{P}_G \triangleq P_G d_G^{-\alpha_s} / (\beta P_S + 1), \\ \bar{P}_{G_s} &\triangleq P_{G_s} d_i^{-\alpha_g}. \end{aligned}$$

When a user requests file W_n , the gateway transfers the non-cached $(1 - \mu_{1,n} - \mu_{2,n})$ portion to the satellite, which will then send accumulated $(1 - \mu_{1,n})$ portion to the ground station. By employing the FastForward (FF) protocol [25] with the FD mode at the satellite, the completion time is $\max\{(1 - \mu_{1,n} - \mu_{2,n})F/R_G, (1 - \mu_{1,n})F/R_S\}$. FF protocol allows the relay node to constructively forward the received data with very low processing latency. An example of the delivery scheme using FF protocol is illustrated in Fig. 3³. Similarly, by deploying the FF protocol with the FD mode at the ground station, the (end-to-end) delivery time for transferring the entire file to the

³We have neglected the propagation delays for simplicity. Otherwise, they can be straightforwardly included in the delivery time expression as approximate constants and the mathematical derivations remain largely unchanged.

Fig. 3: FD transmissions with FastForward relaying.

user is determined by the time for sending the $(1 - \mu_{1,n})$ portion over the satellite downlink to the ground station and the time for sending the whole file over the terrestrial link to the user, i.e.,

$$t_{n,i} = \max \{ F/R_i, (1 - \mu_{1,n})F/R_S, (1 - \mu_{1,n} - \mu_{2,n})F/R_G \}. \quad (7)$$

IV. Successful delivery probability analysis

A. SDP Analysis

The successful delivery probability (SDP) of serving the requested file W_n is the probability that a user receives the file within the user's active time slot. The following analysis is for file W_n and user U_i , and is true for all users.

THEOREM 1. *The closed-form expression for the SDP $\psi_{n,i}$ of file W_n , i.e.,*

$$\psi_{n,i}(\mu_{1,n}, \mu_{2,n}) = \Pr(t_{n,i} \leq T/K) \quad (8)$$

is given in (9).

The probability operator is taken with respect to the channel power gain variables.

Proof:

Substituting (7) into (8) gives rise to

$$\psi_{n,i} = \Pr \left(\max \left\{ \frac{F}{R_i}, \frac{(1 - \mu_{1,n})F}{R_S}, \frac{(1 - \mu_{1,n} - \mu_{2,n})F}{R_G} \right\} \leq \frac{T}{K} \right) = A_1 \times A_2 \times A_3$$

where we have

$$\begin{aligned} A_1 &= \Pr(F/R_i \leq T/K) \\ &= 1 - F_{h_i} \left((2^{\tau/B_{G_s}} - 1) / \bar{P}_{G_s} \right) \\ A_2 &= \Pr((1 - \mu_{1,n})F/R_S \leq T/K) \\ &= 1 - F_{h_S} \left((2^{(1-\mu_{1,n})\tau/B_S} - 1) / \bar{P}_S \right) \\ A_3 &= \Pr((1 - \mu_{1,n} - \mu_{2,n})F/R_G \leq T/K) \\ &= 1 - F_{h_G} \left((2^{(1-\mu_{1,n}-\mu_{2,n})\tau/B_G} - 1) / \bar{P}_G \right) \end{aligned}$$

and $\tau \triangleq F/(T/K)$ is the average file serving rate. ■

We have the following special cases:

- If W_n is cached at ground station, we have

$$\begin{aligned} \psi_{n,i}(1, 0) &= 1 - F_{h_i} \left((2^{\tau/B_{G_s}} - 1) / \bar{P}_{G_s} \right) \\ &= \exp \left((2^{\tau/B_{G_s}} - 1) / (\bar{P}_{G_s} \bar{h}_i) \right). \end{aligned}$$

- If W_n is cached at the satellite, the SDP is

$$\begin{aligned} \psi_{n,i}(0, 1) &= \left(1 - F_{h_i} \left((2^{\tau/B_{G_s}} - 1) / \bar{P}_{G_s} \right) \right) \\ &\quad \times \left(1 - F_{h_S} \left((2^{\tau/B_S} - 1) / \bar{P}_S \right) \right). \end{aligned}$$

- If W_n is fully retrieved from the gateway, the SDP is

$$\begin{aligned} \psi_{n,i}(0, 0) &= \left(1 - F_{h_i} \left((2^{\tau/B_{G_s}} - 1) / \bar{P}_{G_s} \right) \right) \\ &\quad \times \left(1 - F_{h_S} \left((2^{\tau/B_S} - 1) / \bar{P}_S \right) \right) \\ &\quad \times \left(1 - F_{h_G} \left((2^{\tau/B_G} - 1) / \bar{P}_G \right) \right). \end{aligned}$$

The SDP of file W_n under the effect of $\mu_{1,n}$ and $\mu_{2,n}$ is illustrated in Fig. 4. It shows that the SDP can be improved by deploying the second-tier caching at the satellite, and the larger the caching capacity is, the more SDP improvement is. The SDP is the largest or smallest when the file is cached at the ground station or fully retrieved from the gateway, respectively.

Fig. 4: $\psi_{n,i}$ under different caching portions.

The average SDP of the system is defined as a weighted sum of the users' SDPs:

$$\begin{aligned} \psi(\boldsymbol{\mu}_1, \boldsymbol{\mu}_2) &= \sum_{i=1}^K \omega_i \psi_i(\boldsymbol{\mu}_1, \boldsymbol{\mu}_2) \\ &= \sum_{i=1}^K \omega_i \sum_{n=1}^N q_n \psi_{n,i}(\mu_{1,n}, \mu_{2,n}) \end{aligned} \quad (10)$$

with weighting coefficients $\omega_i \in (0, 1)$ and $\sum_i \omega_i = 1$. In the following, without loss of generality, we assume homogeneous users with $\omega_i = 1/K$, and average channel gains $\bar{h}_i = \bar{h}_k, \forall i, k$. Thus, $\psi = \psi_i$.

B. SDP Under Different Caching Policies

In this section, we analyze the average SDP under three different caching policies, i.e., uniform caching,

* SDP of file W_n under generic caching policy:

$$\begin{aligned} \psi_{n,i}(\mu_{1,n}, \mu_{2,n}) &= \exp \left(-(2^{\tau/B_{Gs}} - 1) / (\bar{P}_{Gs} \bar{h}_i) \right) \\ &\times \left(1 - \sum_{k=0}^{m_1-1} \sum_{q=0}^k \frac{(-1)^{2k-q}}{k!q!} \cdot \frac{\alpha_1 \sigma_1^k (1-m_1)_k}{(\sigma_1 - 1/b_1)^{k-q+1}} \cdot \frac{(2^{(1-\mu_{1,n})\tau/B_S} - 1)^{q_e(\sigma_1-1/b_1)} (2^{(1-\mu_{1,n})\tau/B_S} - 1) / \bar{P}_S}{\bar{P}_S^q} \right) \\ &\times \left(1 - \sum_{k=0}^{m_2-1} \sum_{q=0}^k \frac{(-1)^{2k-q}}{k!q!} \cdot \frac{\alpha_2 \sigma_2^k (1-m_2)_k}{(\sigma_2 - 1/b_2)^{k-q+1}} \cdot \frac{(2^{(1-\mu_{1,n}-\mu_{2,n})\tau/B_G} - 1)^{q_e(\sigma_2-1/b_2)} (2^{(1-\mu_{1,n}-\mu_{2,n})\tau/B_G} - 1) / \bar{P}_G}{\bar{P}_G^q} \right) \end{aligned} \quad (9)$$

ground station most popular caching, and satellite most popular caching.

1. Uniform caching policy

The same portions of N files are cached in both tiers. This caching policy is represented by $\mu_1^u = [\mu_1^u, \dots, \mu_1^u]$, $\mu_2^u = [\mu_2^u, \dots, \mu_2^u]$ where $\mu_i^u = N_i/N$. The SDP under this caching policy is

$$\psi^u = \sum_{n=1}^N q_n \psi_{n,i}(N_1/N, N_2/N).$$

The uniform caching policy does not consider the file popularity and tends to be optimal when all files have similar popularity.

2. Ground station most popular caching

The ground station caches N_1 most popular files, and the satellite caches N_2 most popular files of the remaining files. This caching policy considers whole file caching and is represented by $\mu_1^g = [1, \dots, 1, 0, \dots, 0]$, $\mu_2^g = [0, \dots, 0, 1, \dots, 1, 0, \dots, 0]$. The SDP ψ^g is $\sum_{n=1}^{N_1} q_n \psi_{n,i}(1, 0) + \sum_{n=N_1+1}^{N_1+N_2} q_n \psi_{n,i}(0, 1) + \sum_{n=N_1+N_2+1}^N q_n \psi_{n,i}(0, 0)$. This caching policy tends to be better if only a number of files are most popular, i.e., the skewness factor α is large. By caching popular files at the ground station, we can increase the SDP.

3. Satellite most popular caching

The satellite caches N_2 most popular files, and the ground station caches N_1 most popular files of the remaining files. This caching policy also considers whole file caching and is represented by $\mu_1^s = [0, \dots, 0, 1, \dots, 1, 0, \dots, 0]$, $\mu_2^s = [1, \dots, 1, 0, \dots, 0]$. The SDP ψ^s is $\sum_{n=1}^{N_2} q_n \psi_{n,i}(0, 1) + \sum_{n=N_2+1}^{N_2+N_1} q_n \psi_{n,i}(1, 0) + \sum_{n=N_2+N_1+1}^N q_n \psi_{n,i}(0, 0)$. This caching policy is generally more suitable for a two-tier cache-enabled system with multiple ground stations. When most popular files are requested by users associated with different ground stations, the cached files at satellite are broadcast to the ground stations only once. In this case, each ground station does not need to separately cache these popular files. For the system under consideration, we have the following result.

PROPOSITION 1. Assume $N_1 = N_2$. We have $\psi^g > \psi^s$.

Proof:

We have

$$\begin{aligned} \psi^g - \psi^s &= \sum_{n=1}^{N_1} q_n (\psi_{n,i}(1, 0) - \psi_{n,i}(0, 1)) \\ &\quad - \sum_{n=N_1+1}^{N_1+N_2} q_n (\psi_{n,i}(1, 0) - \psi_{n,i}(0, 1)) \\ &= \left(\sum_{n=1}^{N_1} q_n - \sum_{n=N_1+1}^{N_1+N_2} q_n \right) (\psi_{n,i}(1, 0) - \psi_{n,i}(0, 1)) > 0. \end{aligned}$$

■

C. Baselines

1. Single-tier caching

Our derived results can be used to evaluate the SDP of single-tier caching policies. In the single-tier scenario, only cache C_1 is enabled at the ground station. There are two cases:

- When the same portion of N files are cached at the ground station, we have:

$$\psi_s^u = \sum_{n=1}^N q_n \psi_{n,i}(N_1/N, 0).$$

- When N_1 most popular files are cached at the ground station, we have:

$$\psi_s^g = \sum_{n=1}^{N_1} q_n \psi_{n,i}(1, 0) + \sum_{n=N_1+1}^N q_n \psi_{n,i}(0, 0).$$

2. Full-duplex without FastForward relaying

The delivery time of file W_n without the FF protocol is completed in three time slots as in Fig. 5. In the first time slot, the ground station transmits $\mu_{1,n}F$ to the user while receiving $\mu_{2,n}F$ from the satellite; the satellite transmits $\mu_{2,n}F$ to the ground station while receiving $(1 - \mu_{1,n} - \mu_{2,n})F$ from the gateway. In the second time slot, the ground station transmits $\mu_{2,n}F$ to the user while receiving $(1 - \mu_{1,n} - \mu_{2,n})F$ from the satellite. Finally, in the third time slot, the ground station transmits $(1 - \mu_{1,n} - \mu_{2,n})F$

to the user. The delivery time is given by:

$$\begin{aligned} \tilde{t}_{n,i} = & \max \{ \mu_{1,n}F/R_i, \mu_{2,n}F/R_S, (1 - \mu_{1,n} - \mu_{2,n})F/R_G \} \\ & + \max \{ \mu_{2,n}F/R_i, (1 - \mu_{1,n} - \mu_{2,n})F/R_S \} \\ & + (1 - \mu_{1,n} - \mu_{2,n})F/R_i. \end{aligned}$$

Fig. 5: FD transmissions without FastForward relaying.

The SDP of file W_n is $\tilde{\psi}_{n,i} = \Pr(\tilde{t}_{n,i} \leq T/K)$. We will then evaluate the system SDP $\tilde{\psi} = \sum_n q_n \tilde{\psi}_{n,i}$ using Monté Carlo simulation in Section VI.

3. Half-duplex transmission

In this transmission scheme, both the satellite and ground station adopt the HD mode. The achievable rates on the links are

$$\begin{aligned} R_i^{\text{HD}} &= B_{G_s} \log_2 \left(1 + P_{G_s} h_i d_i^{-\alpha_g} \right), \\ R_S^{\text{HD}} &= B_S \log_2 \left(1 + P_S h_S d_S^{-\alpha_s} \right), \\ R_G^{\text{HD}} &= B_G \log_2 \left(1 + P_G h_G d_G^{-\alpha_g} \right). \end{aligned}$$

The timeline for delivering file W_n is divided into three time slots as shown in Fig. 6. The delivery time for file W_n is

$$\begin{aligned} t_{n,i}^{\text{HD}} = & \max \{ \mu_{1,n}F/R_i^{\text{HD}}, (1 - \mu_{1,n} - \mu_{2,n})F/R_G^{\text{HD}} \} \\ & + (1 - \mu_{1,n})F/R_S^{\text{HD}} + (1 - \mu_{1,n})F/R_i^{\text{HD}}. \end{aligned}$$

The SDP of file W_n is $\psi_{n,i}^{\text{HD}} = \Pr(t_{n,i}^{\text{HD}} \leq T/K)$. While $\psi^{\text{HD}} = \sum_n q_n \psi_{n,i}^{\text{HD}}$ is too complex to be found in closed form, it is evaluated using Monté Carlo simulation in Section VI.

Fig. 6: Half-duplex transmission scheme.

D. Energy Consumption of FD and HD Transmission

For a more thorough comparison, we compare the energy consumption of the system in FD/HD modes in addition to the SDP. For the HD transmission scheme, the energies consumption in the three time slots (refer to Fig. 6) are

$$\begin{aligned} E_1^{\text{HD}} &= (P_G/\xi_G + P_{c,G}) (1 - \mu_{1,n} - \mu_{2,n})F/R_G^{\text{HD}} \\ &\quad + (P_{G_s}/\xi_{G_s} + P_{c,G_s}) \mu_{1,n}F/R_i^{\text{HD}}, \\ E_2^{\text{HD}} &= (P_S/\xi_S + P_{c,S}) (1 - \mu_{1,n})F/R_S^{\text{HD}}, \\ E_3^{\text{HD}} &= (P_{G_s}/\xi_{G_s} + P_{c,G_s}) (1 - \mu_{1,n})F/R_i^{\text{HD}}, \end{aligned}$$

with ξ_G, ξ_S , and ξ_{G_s} denote the efficiency of the power amplifier at the gateway, satellite and ground station, respectively; and $P_{c,G}, P_{c,S}$, and P_{c,G_s} are the constant circuit powers corresponding to the satellite uplink, satellite downlink and terrestrial link, respectively. The total energy consumption under the HD mode is

$$E^{\text{HD}} = E_1^{\text{HD}} + E_2^{\text{HD}} + E_3^{\text{HD}}. \quad (11)$$

Different from the HD mode, the FD transmission (with FF relaying protocol) only requires one time slot (refer to Fig. 3) in which all three hops are active. Hence, the total energy consumption is

$$\begin{aligned} E^{\text{FD}} &= (P_G/\xi_G + P'_{c,G}) (1 - \mu_{1,n} - \mu_{2,n})F/R_G \\ &\quad + (P_S/\xi_S + P'_{c,S}) (1 - \mu_{1,n})F/R_S \\ &\quad + (P_{G_s}/\xi_{G_s} + P'_{c,G_s}) F/R_i, \end{aligned} \quad (12)$$

where $P'_{c,j} > P_{c,j}$, $j \in G, S, G_s$ and $(P'_{c,j} - P_{c,j})$ represents the circuit power to implement self-interference cancellation in the FD system.

V. Cache placement design

CACHED contents are placed into caches C_1 and C_2 via satellite uplink and downlink transmissions during off-peak hours. In this section, the cached portions $\mu_{1,n}$ and $\mu_{2,n}$ are designed to maximize the SDP. The cache

placement design problem subject to caching capacity constraints is formulated as follows:

$$\max_{\mu_1, \mu_2} \psi(\mu_1, \mu_2) \quad \text{s.t.} \quad \sum_{n=1}^N \mu_{i,n} \leq N_i, i = 1, 2. \quad (13)$$

(13) is a non-linear non-convex optimization, and hence, it is very challenging to find the optimal solution, which can be found using exhaustive search only in general. In the following, we resort to simpler caching placement solutions assuming that the cached files are prefetched with the same portions in each tier. We consider the following two policies.

Generalized uniform caching policy. The caching policy can be described by the vectors $\bar{\mu}_i^u = [\bar{\mu}_i^u, \dots, \bar{\mu}_i^u, 0, \dots, 0]$, where $\bar{N} \leq N$ and $\bar{\mu}_i^u = N_i/\bar{N}$, $i =$

1, 2. In this policy, $\bar{N}$ most popular files are being cached with the same portions at the ground station and the satellite. This generalized uniform caching policy is inspired by the fact that uniform caching policy (Section 1) naively caches all files in the library regardless of the file request rate. The cache resource is wasted for prefetching files that are rarely requested. Comparing to the uniform caching policy, this generalized uniform caching policy prioritizes more popular files. This caching policy becomes the uniform caching policy when $\bar{N} = N$.

Generalized ground station most popular caching policy. The same $\bar{\mu}_1^g$ portion of $\bar{N}_1$ most popular files are cached at ground station, and $\bar{\mu}_2^g$ of $\bar{N}_2$ most popular files of the remaining files are cached at satellite. This policy can be represented by $\bar{\mu}_1^g = [\bar{\mu}_1^g, \dots, \bar{\mu}_1^g, 0, \dots, 0]$;

and $\bar{\mu}_2^g = [0, \dots, 0, \bar{\mu}_2^g, \dots, \bar{\mu}_2^g, 0, \dots, 0]$, $\bar{\mu}_i^g = N_i/\bar{N}_i$.

Comparing to the whole file caching in the ground station most popular policy (Section 2), the content cached in this policy is partial-file, allowing more popular files being prefetched in the two cache tiers. Hence, the generalized ground station caching policy is more diverse in terms of file popularity. This caching policy boils down to the ground station most popular caching policy when $\bar{N}_i = N_i$, $i = 1, 2$.

We can conduct a simple exhaustive search to find the generalized policy to maximize the SDP in (13), i.e., by searching the number of files $\bar{N}$ or $\bar{N}_i$ to be pre-fetched. It can be seen that the obtained policy allows more popular files to be cached closer to users using limited caching resources. As a result, the serving time can be reduced.

REMARK 1. *The obtained caching solution always achieves a higher SDP than the ground station most popular and uniform caching strategies.*

VI. Numerical results

THE values for key parameters are presented in Table I for illustrative purpose [26]. As more than 100 dB of self-interference (SI) cancellation level is required for the

FD system to achieve similar signal-to-noise-ratio-plus-interference as that of the HD system [27], SI cancellation parameter β is chosen to be 10^{-4} to achieve SI cancellation level about 110 dB to 120 dB in our FD system setting. The fading states for the satellite-terrestrial links are defined by $\{m_1, b_1, \Omega_1\} = \{5, 0.251, 0.279\}$ [20] and $\{m_2, b_2, \Omega_2\} = \{4, 0.126, 0.835\}$ [21] when approaching average shadowing. End-users are randomly distributed inside a circle with radius 10 km and center at G_s .

TABLE I: Parameters Used for Numerical Results

Parameter	Value	Parameter	Value
N	1000 files	F	100 Mb
$B_G/B_S/B_{G_s}$	2/10/2 GHz	$P_S/P_G/P_{G_s}$	10/30/2 W
σ^2	-120 dBm/Hz	d_S/d_G	35,786 km
α	0.8	α_s/α_g	2/3
$m_1/b_1/\Omega_1$	5/0.251/0.279	$m_2/b_2/\Omega_2$	4/0.126/0.835
ξ_i	0.35	$P_{c,i}/P'_{c,i}$	2.5/3 W

A. FD versus HD: Effect of SI Cancellation Parameter β

We first investigate the average SDP when β varies. Consider $C_1 = 50$ Gb, $C_2 = 10$ Gb, and $K = 20$ users under uniform caching policy.

Fig. 7: Effect of the SIC parameter on average SDP.

With the HD scheme as a reference, the effect of β on the FD scheme SDP is displayed in Fig. 7. As β increases (i.e., worse self-interference cancellation in the FD system), the SDP decreases and starts getting worse than that of HD transmission. The FD scheme achieves the same SDP as that of HD with $\beta = 1.7 \times 10^{-3}$. The self-interference cancellation in the FD system in this case is 103 dB to 107 dB which is almost the recommended value in [27]. The SDP of FD is 50% worse than that of HD with $\beta = 10^{-1}$ where self-interference cancellation is 85 dB to 90 dB. We also present numerical result for hybrid FD/HD transmission. In this hybrid scheme, the channel realizations are generated, then the delivery

time of each file is the minimum delivery time between the FD and HD modes. In practical applications, once knowing the channel condition, devices in the system can dynamically switch between the FD and HD modes for better performance.

Fig. 8: Effect of the SIC parameter on energy consumption.

Fig. 8 shows the system total energy consumption when β varies. With $\beta = 10^{-4}$, which indicates to good SI cancellation, the two systems consume almost the same amount of energy but the FD system achieves 38% higher in SDP. When β increases, the energy consumption in FD also increases and the SDP decreases. This amount of energy accounts for the SIC process and the energy for transmitting unsuccessful files.

B. Two-tier versus Single-tier Caching

The single-tier cache-enabled system in this section is adopted from STN system models in [8]–[10], where content caches are placed at ground stations only. The caching policies for the single-tier cache system correspond to those at the ground station in the two-tier cache system. For a fair comparison between the two-tier and single-tier caching systems, we set up both systems with the same total cache capacity of 80 Gb serving 20 users. The average SDPs of the uniform and ground station most popular caching policies are shown in Fig. 9.

Under the uniform caching policy, N_1/N varies and N_2/N is set to be 0.2 in the two-tier caching (meaning the satellite cache capacity is 20 Gb). The result is shown in Fig. 9 with the two-tier scheme outperforming its single-tier counterpart. At best performance with $N_1/N = 0.7$, the achieved SDP in the two-tier scheme is 25% higher than that of the single-tier scheme. Under the ground station most popular caching policy, the number of most popular file cached N_1 varies; N_2 is fixed at 200 files in the two-tier scheme. The two-tier scheme gives approximately 40% SDP higher than that of the single-tier scheme in best performance scenarios with the ground station caching 600 most popular files. The single-tier system only outperforms the two-tier system when the

Fig. 9: Total caching capacity comparison.

Fig. 10: Total energy consumption comparison.

total caching capacity is 200 files under uniform caching policy. In this case, we are comparing single-tier system with content cached solely at ground station and two-tier system with content cached at satellite only. The corresponding energy consumption is shown in Fig. 10. With more files cached, both systems consume less energy. Under ground station most popular caching policy, both systems achieve higher SDP while consume less energy than under uniform caching policy.

C. FastForward versus without FastForward Relaying

To compare the performance of the FD protocols, the system is set up with varying values for C_1 and C_2 of 20 Gb serving 20 users. Full-duplex with and without the FastForward relaying protocols are examined under the uniform and ground station most popular caching policies. File popularity is set to be α of 0.8 and β of 10^{-4} for the FD mode. As shown in Fig. 11, the system with the FastForward relaying protocol achieves a higher SDP. This result is expected since the delivery time using the FastForward protocol is shorter than that without using FastForward relaying.

Fig. 11: Average SDP under FD with and without the FastForward relaying protocols.

D. Network Capacity

To examine the serving capability, the system is set up with the two layers' caches C_1 of 50 Gb and C_2 of 20 Gb. The number of users K varies while other parameters are listed in Table I. Fig. 12 shows the number of supportable users by the two-tier system under two caching policies. A smaller average SDP under uniform caching is achieved when supporting a small number of users; however, it surpasses the SDP of the ground station most popular caching scheme with a larger number of users. The acceptable SDP value depends on practical applications. In this case, with SDP of 50%, the results follow a common understanding that the uniform caching scheme is not as effective as the ground station most popular caching scheme.

Fig. 12: Number of supportable users under different caching policies.

E. Caching Capacity

In this section, we analyze the two-tier caching capacity when serving 20 users. The system is configured using the parameters specified in Table I. The cache capacities

of the two layers are shown in Fig. 13. As expected, the ground station most popular caching policy outperforms uniform caching. The more cache capacity is allocated to the ground station, the higher the SDPs are attainable; and the more content is cached at the satellite, the more the SDPs improve. Under uniform caching, the highest SDP is achieved at 70%, while it is 80% under ground station most popular caching.

Fig. 13: Effect of caching capacity on the average SDP.

F. Effect of File Popularity

In this subsection, to examine the effect of the file popularity on the SDP under the ground station and satellite most popular caching policies, the system is configured with two tier caches C_1 and C_2 of 50 Gb and 20 Gb, respectively, and serves 20 ground users. Through varying Zipf skewness factor α , the resulting SDP is plotted in Fig. 14. Higher SDPs are achieved for both most popular caching schemes with an increase in the Zipf skewness factor. The results also match Proposition 1. This is because having one ground station in the system, the ground station most popular caching policy always gives better SDP results than under satellite most popular caching.

Fig. 14: Effect of file popularity on the average SDP.

G. Effect of the File Size

To investigate the effect of the file size on SDP, we vary the sizes of N files at the gateway following a uniform distribution with a mean of 100 Mb and a standard deviation of 50 Mb. Two caching schemes are set up as follows: (i) C_1 is prefetched with the largest files and then C_2 is filled with the largest files among the remained files; (ii) C_2 is prefetched with the largest files and then C_1 is filled with the largest files among the remained files. The results are shown in Fig. 15. The average SDP is better when the largest files are cached at the ground station, resulting in a gain of approximately 10% to 20% in SDP.

Fig. 15: Effect of the file size on the average SDP.

H. Generalized Caching Solutions

To investigate the performances of the generalized caching policies, we find the solutions when the caching capacity at satellite is fixed at 20 Gb and at ground station varies up to 80 Gb. The performances are illustrated in Fig. 16. As can be seen from the figure, generalized caching policies yield higher SDPs than the naive uniform and ground station most popular caching policies. The solutions for each case of C_1 is detailed in Table II.

TABLE II: Solutions of The Generalized Caching Policies.

C_1 (Gb)	Generalized uniform				Gen. grnd. sta. most pop.			
	N_1	$\bar{\mu}_1$	N_2	$\bar{\mu}_2$	N_1	$\bar{\mu}_1$	N_2	$\bar{\mu}_2$
10	250	0.40	400	0.50	220	0.45	370	0.54
20	350	0.57	500	0.40	300	0.67	370	0.54
30	450	0.66	650	0.30	500	0.60	330	0.60
40	650	0.61	650	0.30	570	0.70	300	0.66
50	700	0.71	800	0.25	600	0.83	280	0.71
60	800	0.75	900	0.22	600	1.00	200	1.00
70	910	0.76	900	0.22	700	1.00	200	1.00
80	950	0.84	1000	0.15	800	1.00	200	1.00

Fig. 16: Average SDP under the generalized caching policies.

I. Optimal Caching Solution

Under generic caching policy, the original optimization problem in (13) could be solved in case of a small number of files N . For the purpose of illustration, we compute the optimal solution for the problem in (13) in the case of $N = 20$ in Fig. 17. The caching capacity at satellite is fixed at 1 Gb and at ground station varies up to 4 Gb. The optimal solution (generic caching) of the original problem in (13) is found using exhaustive search. For the optimal solution, with one constraint on the capacity of the two cache tiers C_1 and C_2 , we generate feasible caching solutions by computing its SDP values using (9), and then determining the best one.

Fig. 17: Optimal caching solution.

As expected, the optimal solution under generic caching gives the best SDP comparing to the sub-optimal solution under the two generalized caching policies. The results show that the gaps between the optimal and sub-optimal solutions are quite negligible, demonstrating the effectiveness of the caching policies under consideration at significantly reduced computational complexity.

VII. Conclusion

IN this paper, we studied the SDP of a two-tier cache enabled satellite IoT system with full-duplex communications. A thorough analysis for the SDP under the FD mode with the FastForward relaying protocol was carried out under different caching policies. Cache placement was also designed with the goal of maximizing the average SDP. Using the FD mode without the FastForward protocol, HD mode and single-tier cache system as benchmarks, numerical results were presented when investigating the system performance in terms of SDP, maximum number of supportable users, caching capacity in each tier, and the effect of file popularity on the system performance under different caching policies. For a better understanding on parameters that affect the proposed system performance, a special case with a variety of file sizes was also investigated.

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